

# NYU Poly Reverse Engineering Lecture Session II

Aaron Portnoy

TippingPoint Security Research

Peter Silberman

Mandiant Engineering and Research

# Intro to x86

- Contains 8 general purposes registers
  - @eax, ebx, ecx, edx, esi, edi, ebp, esp
  - Consider them temporary variables
- Stack is used to store registers when values need to be saved
  - LIFO (push/pop)

# Basic Program Execution Registers

## Eight General Purpose Registers

EAX	EBP
EBX	ESP
ECX	ESI
EDX	EDI

## Processor Status Flags

EFLAGS
--------

## Instruction Pointer

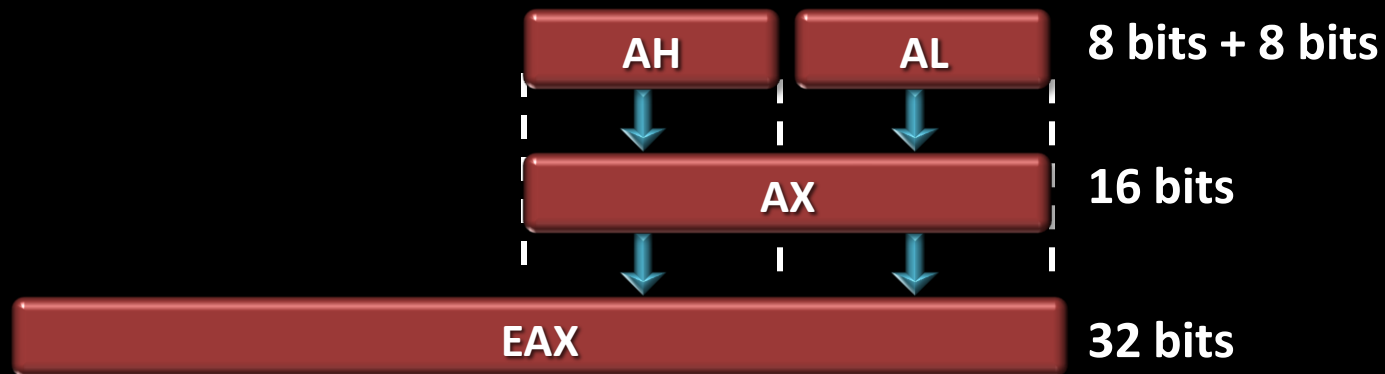
EIP
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## Six Segment Registers

CS	ES
SS	FS
DS	GS

# General Purpose Registers

- The 8 general purpose registers are used for arithmetic and data movement
- Each register can be addressed as a 32 bit, 16 bit or 8 bit value



# Overlapping Registers

32-bit	16-bit	8-bit (high)	8-bit (low)
EAX	AX	AH	AL
EBX	BX	BH	BL
ECX	CX	CH	CL
EDX	DX	DH	DL
ESI	SI		
EDI	DI		
EBP	BP		
ESP	SP		

# General-Purpose Registers

- Some of these registers are used by specific instructions
  - EAX is automatically used by multiplication and division operations
  - ECX is used as a counter in several instructions
  - ESI and EDI are as src and dst for copying data in loops
  - EBP and ESP are used to track changes to the stack
- Calling conventions and ABIs define certain registers uses
  - EAX is used to store the return value for function calls
  - ECX is used to store a pointer to the 'this' object in C++

# Classes of Instructions

- X86 has a lot of instructions
  - We are only going to cover a select few
- Instructions that:
  - Read
  - Write
  - Compare
  - Branch
  - Perform Arithmetic
    - Add
    - Subtract
    - Multiply
    - Divide
    - Bitmath
    - Floating Point

# Read Instructions

- Memory dereferencing is the equivalent of reading

unsigned long x = 0;

x = \*p;

- Where \* is dereference assignment
- [ ] is dereference assignment in x86

mov reg32, [reg32] -> mov eax, [ecx]

mov reg32, [imm32] -> mov eax, [04010000]

pop eax



# Write Instructions

- Consider these a store
  - `mov [ebx], 0x20`
    - Stores the immediate value 0x20 at the address specified by `@ebx`
  - `mov [ecx+0x14], edx`
    - Stores the value of the `edx` register into the address at `ecx` plus 0x14 bytes

# Stack Operations

- push
  - Syntax: *push src*
- Examples:
  - push eax
  - push 0x100
  - push dword\_0x100400
- pop
  - Syntax: *pop dst*
- Examples:
  - pop eax

# Arithmetic

- `inc`
  - Syntax: *inc dst*
- Examples:
  - `inc edx`
- `dec`
  - Syntax: *dec dst*
- Examples:
  - `dec eax`

# Arithmetic (cont.)

- `mul`
  - Syntax: *mul src*
    - *Result is stored in ecx*
- Examples:
  - `mul edx`
- `div`
  - Syntax: *div src*
    - *Result is stored in eax*
- Examples:
  - `div edi`

# Arithmetic (cont.)

- add
  - Syntax: *add dst, src*
- Examples:
  - add eax, 10
  - add edx, eax
- sub
  - Syntax: *sub dst, src*
- Examples:
  - sub eax, 10
  - sub ecx, edx

# Arithmetic (cont.)

- `lea`
  - Syntax: *lea dst, src*
- Examples:
  - `lea eax, [eax*4]`
  - `lea edx, [edi+ecx]`

# Arithmetic (cont.)

- Bitmath

- shl, shr

- Shifts the bits of the operand either to the left or to the right

- $00000111 \ll 2 = 00011100$

- Examples:

- shl eax, 2
    - shr edx, 4

# Arithmetic (cont.)

- Floating Point
  - *fbld, fild, fcmovnb, ...*
  - Consult x86 manuals



# Comparisons

- `cmp`
  - Syntax: *cmp dst, src*
- Examples:
  - `cmp eax, ecx`
  - `cmp edx, 10`
- `test/and`
  - Syntax: *test dst, src*
- Examples:
  - `and eax, 10`
  - `test ecx, edx`

# Branches

- Used to direct code, frequently based on previous comparison
- `jxx`
  - Syntax: *jxx dst*
- Examples:
  - `jz reg32`
  - `jnb $-5`
  - `jnz 0x04010012`

# Code Execution Transfers

- `Call`
  - Used to call functions
  - Syntax: *call src*
- Examples:
  - `call ecx`
  - `call [40100000]`
  - `call 0x41c2200c`
- Occasionally the `jmp` instruction will be responsible for transferring execution to another function
  - Examples:
    - `jmp 0x41c2200c`

Questions?

# Vulnerability Classes

- Those we'll cover
  - Integer Overflows
  - Stack/Heap Based Buffer Overflow
  - Format Strings
- Those we won't
  - Invalid Free/Double Free
  - Uninitialized Variables
  - Misc. (memory corruption)
    - MS09-028: <http://bit.ly/owningMSdirectshow>

# Integer Wraps

- Integers are able to store a finite size
- Integer wraps due to type conversion
  - Width
    - Unsigned long to short
- Integer wraps due to arithmetic
  - $\text{MAX\_INT} + x$
  - $0 - x$

# Integer Overflows (Ex 1)

```
unsigned long a = 0xFFFFFFFFFD;  
unsigned long b = 3;  
unsigned long c = 2;  
c = a + b;  
printf("Result: 0x%08x\n", c);
```

- Output:  
*Result: 0x00000000*

# Integer Overflows (Ex1)

```
var_c      = dword ptr -0Ch
var_b      = dword ptr -8
var_a      = dword ptr -4

push       ebp
mov        ebp, esp
sub        esp, 0Ch
mov        [ebp+var_a], 0FFFFFFDh
mov        [ebp+var_b], 3
mov        [ebp+var_c], 2
mov        eax, [ebp+var_a]
add        eax, [ebp+var_b] ; a + b
mov        [ebp+var_c], eax ; c = (result of a + b)
mov        ecx, [ebp+var_c] |
push       ecx
push       offset format    ; "Result: 0x%08x\n"
call       _printf
add        esp, 8
mov        esp, ebp
pop        ebp
retn
```



# Integer Overflows (Ex2)

```
void SomeFunc(unsigned long user_supplied, char * userbuffer)
    unsigned long a = 0;
    SHORT b = 0;
    char mybuffer[300];

    a = user_supplied;
    b = a;
    printf("Checking %d to make sure it is less than 300\n", b);
    if(b >= 300)
    {
        printf("Thank you come again..\n");
        return;
    }

    printf("Passed my checks, copying %d bytes into buffer of size %d\n", a, sizeof(mybuffer));
    strncpy(mybuffer, userbuffer, a);
```

# Integer Overflows (Ex2)

- Output:

Checking 0 to make sure it is less  
than 300

Passed my checks, copying 65536 bytes  
into buffer of size 300

# Stack/Heap Overflows

- Stack/Heap overflows are the most common memory mismanagement
  - Smashing the Stack for Fun and Profit, classic:  
<http://www.phrack.org/issues.html?id=14&issue=49>
  - Gera's insecure programming examples:  
<http://community.corest.com/~gera/InsecureProgramming/abo1.html>
- Essentially, exceeding the bounds of an allocation during a sequence of writes

# Stack/Heap Overflows (Ex1)

```
dst      = byte ptr -68h
src      = dword ptr  8

push     ebp
mov      ebp, esp
sub      esp, 68h
mov      eax, [ebp+src]
push     eax                ; src
lea      ecx, [ebp+dst]
push     ecx                ; dst
call     _strcpy
add      esp, 8
mov      esp, ebp
pop      ebp
retn
```

# Stack/Heap Overflows (Ex2)

```
.text:10001341 mov     ebx, dword_10001344
.text:10001347 push    ebp
.text:10001348 mov     ebp, [eax+8]
.text:10001349 mov     eax, [eax+4]
.text:1000134E push    esi
.text:1000134F push    edi
.text:10001350 mov     edi, [eax+10h]
.text:10001353 push    8 ; Size
.text:10001355 call    ebx ; malloc
.text:10001357 lea     ecx, [edi+edi]
.text:1000135A movsx   edx, cx
.text:1000135D inc     edx
.text:1000135E push    edx ; Size
.text:1000135F mov     esi, eax
.text:10001361 call    ebx ; malloc
.text:10001363 add     esp, 8
.text:10001366 test    eax, eax
.text:10001368 jz      short loc_1000138C
.text:1000136A push    edi ; Size
.text:1000136B push    ebp ; Src
.text:1000136C push    eax ; Dst
.text:1000136D mov     [esi], eax
.text:1000136F call    memcpy
.text:10001374 mov     eax, dword_10001344
.text:10001379 add     esp, 0Ch
.text:1000137C mov     [esi+4], eax
```

# Format String

- Passing unsanitized user input to a function that accepts a format string
  - Caused by C's ability to use varargs
  - Enables attackers to read or write data to memory
  - %s, %x, %n

# Format String (Ex1)

```
void pretty_print(char * user)  
    printf(user);
```

- No sanitization of user input what happens we direct the function to continually pop items off the stack?

Questions?



Automation

# Vulnerability Hunting Styles

- *I've been up for 3 days straight, where's my coffee: Cerebral and Successful Method (Aaron)*
  - Fully reverse document a product's internal workings
  - Pros:
    - Full understanding of the product
    - In the developers head
    - Finding bugs is much easier, think like the developer
  - Cons:
    - Time consuming

# Vulnerability Hunting Styles

- *Where's my Ritalin: ADHD Induced Method, Less successful (Peter):*
  - Only reverse points of input:
    - CreateFile, recv, ReadFile, rpc, etc..
  - Spend sometime understanding how to craft input to get most code coverage
  - Pros:
    - Finds bugs faster
    - less time spent reversing
    - Don't have to consume A LOT of coffee
  - Cons:
    - Only knocks off low to medium hanging fruit.
    - May not get full code coverage

# Vulnerability Hunting Styles

- As much as our styles differ, we both make use of automation
  - Using conditional breakpoints to gather information
  - Analyzing binary code programmatically
  - Instrumenting an application
    - Call `unknown()` 10000 times with differing args, analyze output
  - Gathering runtime data
    - Fill in cross references with dynamic call information
    - Dump global variable values

# Automating Binary Analysis

- Binary analysis tasks automated for speed
  - Deobfuscation
  - Control and Data Flow Analysis
    - CF Analysis can lead to identification or programming errors
      - Bad calls, dangerous loops, signed/unsigned compares etc..
    - DF Analysis can be used for type reconstruction or to improve CF analysis
      - User supplied variable + 2 = integer overflow
        - » Knowing its user supplied is the key

# Automating Binary Analysis

- Pattern Matching:
  - Can be used to auto comment commonly used instruction sequences
    - Inline strcpy/strcmp/strlen
    - Inline memcpy, memset, memmove
    - FindCrypt plugin (for locating common crypto methods)
  - Find possibly interesting code to audit
    - Arithmetic followed by allocations
    - Format string calls with no format token

# IDA + Automated Binary Analysis

- IDA provides its own scripting interface called IDC
- IDAPython gives users access to the SDK and IDC
  - All in python!
  - This is what we will focus on today
- Note we will be using IDAPython 2.5 (compiled from trunk)
  - <http://thunkers.net/~deft/misc/Reversing102.zip>

# Iterating Over Functions

```
for ea_start in Functions(MinEA(), MaxEA()):  
    print "%s: 0x%08x" % (GetFunctionName(ea_start), ea_start)
```

- **Output:**

```
__SEH_epilog: 0x757319fc  
__SEH_prolog: 0x75731a0d  
?LsapScavengerTrigger@@YGKPAX@Z: 0x75731a4d  
?LsapTimerCallback@@YGXPAXE@Z: 0x75731ac1  
?LsapDerefScavItem@@YGXPAU_LSAP_SCAVENGER_ITEM@@@Z: 0x75731b46  
?LsapScavengerBreak@@YGKPAX@Z: 0x75731b77  
_SafeAllocFreeToHeap@4: 0x75733086  
_LsapAllocateLsaHeap@4: 0x757330ab  
(..)
```



# Iterating Over Function's Basic Blocks

- Still a little bit buggy
- Only enabled in IDAPython pulled from trunk

# Iterating Over Function's Basic Blocks

```
func = get_func(get_screen_ea())
fc = idaapi.FlowChart(func)

for block in fc:
    print "%X - %X [%d]:" % (block.startEA, block.endEA, block.id)

    for succ_block in block.succs():
        print "    %X - %X [%d]:" % (succ_block.startEA,
succ_block.endEA, succ_block.id)

    for pred_block in block.preds():
        print "    %X - %X [%d]:" % (pred_block.startEA,
pred_block.endEA, pred_block.id)
```

# Iterating Over Function's Instructions

```
func = get_func(get_screen_ea())
for ea in Heads(func.startEA, func.endEA):
    print "0x%08x: %s" % (ea, GetDisasm(ea))
```

- **Output:**

```
0x75789146: mov     edi, edi
0x75789148: push    ebp
0x75789149: mov     ebp, esp
0x7578914b: lea     eax, [ebp+arg_0]
0x7578914e: push    eax
0x7578914f: call    _LsapUnregisterAuditEvent@4
0x75789154: pop     ebp
0x75789155: retn    4
```

# Tying it all together

- Bad call scanner (I know sooo 2000 and late)
- Identifies dangerous calls to known bad APIs  
i.e., strcpy, sprintf etc.

# Tying it all together

- Finding PE Parsing routines:

```
for start in Functions(MinEA(), MaxEA()):
    for ea in Heads(start, PrevAddr(get_func(start).endEA)):
        disasm = GetDisasm(ea)
        name = GetFunctionName(ea)

        if disasm.lower().find("5a4d") != -1:
            l = "%s => 0x%08x: %s\n" % (name, ea, disasm)
            msg(l)

        if disasm.lower().find("4550") != -1:
            l = "%s => 0x%08x: %s\n" % (name, ea, disasm)
            msg(l)
```

# Tying it all together

- Reverser's Cookbook methods
  - find\_path, find\_all\_paths
  - find\_instr, find\_func
  - enum\_switches
  - file\_io, net\_io
  - ...

Questions?

# Debugging



# Debugging Crashes

- When a crash occurs we are mostly concerned with...
  - Faulting instruction
    - Disassembly around faulting instruction (ub @eip in windbg)
  - Register contents
    - Pointers
      - To code
      - To data
    - Values
      - Simple data types
      - Return values
      - Lengths/Counters
  - Call stack
    - How we got here
  - State of the heap and stack
    - Verify heap integrity (!heap 0 -v)
    - Verify stack integrity (dd @esp ; !exchain)

# Debugging Crashes (cont.)

- VM Debugging
  - Allows you to snapshot an entire system
    - This allows you to know ahead of time where things are allocated
  - Revert back to a good state

# Clever Breakpoints

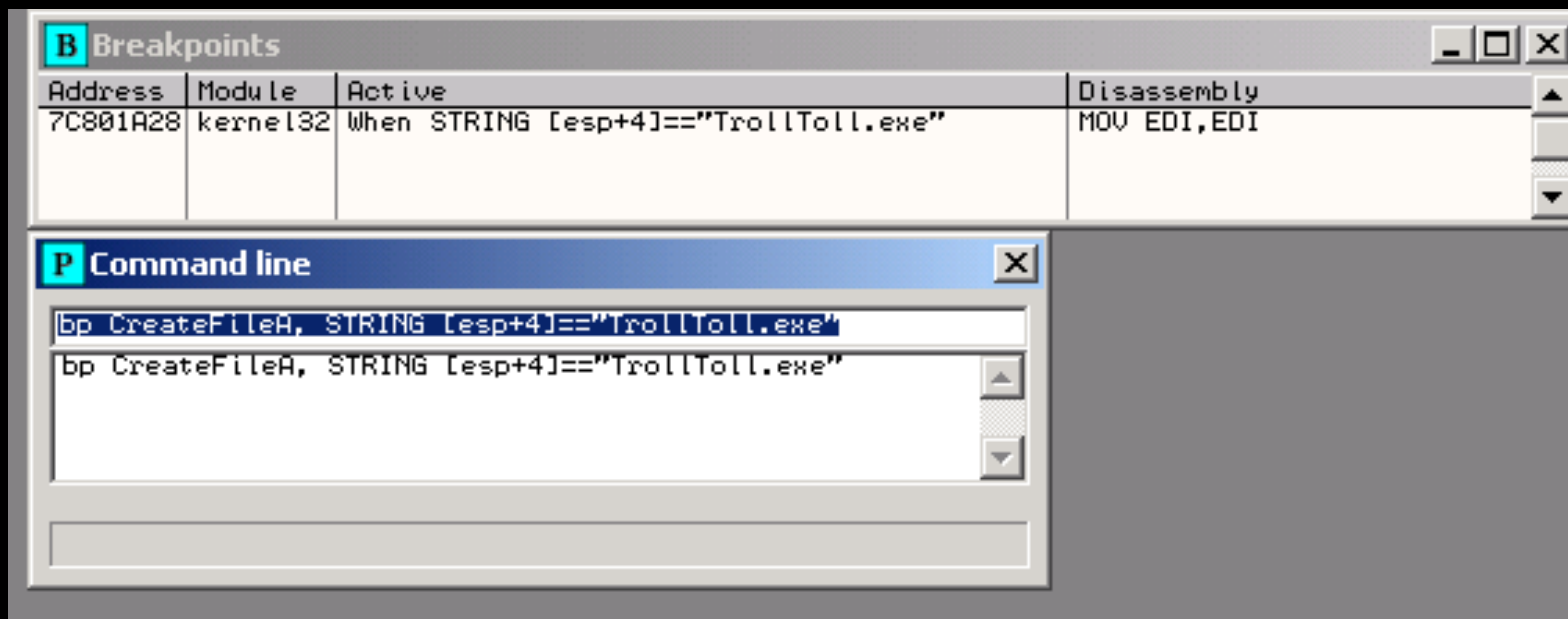
- Breakpoints are great for examining a program
  - Can get an inside peak at what is going on
    - Tracing! (check your log window)
  - Break if loop counter == 0xFFFD
- Memory breakpoints
  - When is a buffer first read from? Written to?
  - Useful for tricks like hunting down the source of an allocation
    - VM debugging helps here

# Conditional Breakpoints (WinDbg)

- Conditional Breakpoints
  - Breakpoint is only executed when condition is met:
    - `bp myprogram!SomeFunc+0x08 "j @eax = 0xFFFFD ';'gc"`
      - <http://msdn.microsoft.com/en-us/library/cc267482.aspx>

# Conditional Breakpoints (OllyDbg)

- OllyDbg has a command line plugin
  - Bp CreateFileA, STRING [esp+4] == "TrollToll.exe"



# Conditional Breakpoints (GDB)

- GDB standard \*nix debugger (unfortunately)
  - b \*0x00401000 if \$eax==5

# Useful Debugging Tricks

- GFLAGS
  - Page Heap (awesome)
    - Can obscure vulnerabilities, though
      - Pointer math
    - Enables the debug heap, behaves differently
  - User mode stack trace database
    - Track sources of heap allocations
  - Heap \* checking
    - Ensures heap integrity during heap operations
  - Read up: [http://technet.microsoft.com/en-us/library/cc738763\(WS.10\).aspx](http://technet.microsoft.com/en-us/library/cc738763(WS.10).aspx)

# Useful Debugging Tricks

- WinDBG
  - !heap
    - Walk the heap (requires symbols)
    - !heap 0 -v
    - !heap -p -a 0xwhatever
  - !exchain
    - Lists the registered exception handles
  - !analyze -v
    - Analyzes the current crash, includes call stack, other useful info
  - !exploitable
    - Not as cool as it sounds
- ImmDbg
  - Has it's own !heap



Questions?

# Hands On Experience: GreenMan



# Background

- GreenMan
  - Compiled for Linux/Mac/Windows
  - Vulnerable Application (sorry to kill the suspense)
  - Listens on port 4959
  - GreenMan has the following vulnerabilities:
    - Integer overflow resulting in a heap overflow
    - Stack overflow
    - Format string

# Background

- SweetDee.py
  - Client for GreenMan
  - `sweetdee.py <host> <port> <opcode> <payload>`
    - Host – your ip OR ALL which sends a packet to the ip range:
      - 192.168.1.0-255
    - Port – should be 4959 unless recompiled
    - Opcode – which vulnerability do you want to trigger?
    - Payload (optional) – used only in opcode 3

# You need to...

- Download
  - <http://thunkers.net/~deft/misc/Reversing102.zip>
- Please connect to Wireless Access Point
  - VirusNetwork
- Disable your firewalls ☺
- Run the GreenMan application
- Attach your debugger of choice to GreenMan
  - WinDBG: F6
  - OllyDBG:File->Attach
  - GDB: gdb `pidof GreenMan`

# Reversing the Binary

- Open up the GreenMan binary in IDA
  - Windows/Linux/OS X versions will look different
- Check the Imports
  - This will give you an idea of what functions it uses
  - However, if it's statically compiled (linux) it won't have Imports
    - So check Names subview
  - Check exports to find main(). Double-click it.

# Reversing the Binary (Linux)

```
.text:08048426      public main
.text:08048426 main      proc near                ; DATA XREF: _start+17↑to
.text:08048426      fd          = dword ptr -2760h
.text:08048426      addr         = dword ptr -275Ch
.text:08048426      len          = dword ptr -2758h
.text:08048426      arg_0        = byte ptr 4
.text:08048426
.text:08048426      lea         ecx, [esp+arg_0]
.text:0804842A      and         esp, 0FFFFFFF0h
.text:0804842D      push        dword ptr [ecx-4]
.text:08048430      push        ebp
.text:08048431      mov         ebp, esp
.text:08048433      push        ecx
.text:08048434      sub         esp, 2754h
.text:0804843A      mov         [esp+2760h+len], 0
.text:08048442      mov         [esp+2760h+addr], 1
.text:0804844A      mov         [esp+2760h+fd], 2
.text:08048451      call        socket
.text:08048456      mov         [ebp-0Ch], eax
.text:08048459      cmp         dword ptr [ebp-0Ch], 0FFFFFFFFh
.text:0804845D      jnz         short loc_8048477
.text:0804845F      mov         [esp+2760h+fd], offset aSocketFailed ; "socket() failed"
.text:08048466      call        puts
.text:0804846B      mov         [esp+2760h+fd], 1
.text:08048472      call        exit
.text:08048477 ; -----
.text:08048477
.text:08048477 loc_8048477:                ; CODE XREF: main+37↑j
.text:08048477      mov         [esp+2760h+len], 10h
.text:0804847F      mov         [esp+2760h+addr], 0
.text:08048487      lea         eax, [ebp-2734h]
.text:0804848D      mov         [esp+2760h+fd], eax
.text:08048490      call        memset
.text:08048495      mov         word ptr [ebp-2734h], 2
.text:0804849E      mov         dword ptr [ebp-2730h], 0
.text:080484A8      mov         [esp+2760h+fd], 4959
.text:080484AF      call        ntohs
.text:080484B4      mov         [ebp-2732h], ax
.text:080484BB      lea         eax, [ebp-2734h]
```

# Reversing the Binary (Linux)

- Things to note
  - Rather than push, gcc will move the addresses directly to the address of esp+X

```

; Attributes: bp-based frame

public DoBufferOverflow
DoBufferOverflow proc near

var_A= byte ptr -0Ah
arg_0= dword ptr  8

push    ebp
mov     ebp, esp
sub     esp, 18h
mov     eax, [ebp+arg_0]
add     eax, 4
mov     [esp+4], eax
lea     eax, [ebp+var_A]
mov     [esp], eax
call    strcpy
mov     eax, 0
leave
retn
DoBufferOverflow endp
```



# Reversing the Binary (Linux)

- Things to note
  - Rather than push, gcc will move the addresses directly to the address of esp+X

```
; Attributes: bp-based frame

public DoBufferOverflow
DoBufferOverflow proc near

var_A= byte ptr -0Ah
arg_0= dword ptr  8

push    ebp
mov     ebp, esp
sub     esp, 18h
mov     eax, [ebp+arg_0]
add     eax, 4
mov     [esp+4], eax
lea     eax, [ebp+var_A]
mov     [esp], eax
call    strcpy
mov     eax, 0
leave
retn
DoBufferOverflow endp
```

# Reversing the Binary (Linux)

- Things to note
  - Even though source code uses `recv()` and `printf()` the compiler substitutes

# Reversing the Binary (Linux)

```
.text:08048516
.text:08048516 loc_8048516:                                ; CODE XREF: main+180↓j
.text:08048516                                           ; main+1CB↓j ...
.text:08048516      lea     eax, [ebp-2724h]
.text:0804851C      lea     edx, [ebp-2744h]
.text:08048522      mov     [esp+2760h+len], eax
.text:08048526      mov     [esp+2760h+addr], edx
.text:0804852A      mov     eax, [ebp-0Ch]
.text:0804852D      mov     [esp+2760h+fd], eax
.text:08048530      call    accept
.text:08048535      mov     [ebp-8], eax
.text:08048538      cmp     dword ptr [ebp-8], 1
.text:0804853C      jnz     short loc_8048556
.text:0804853E      mov     [esp+2760h+fd], offset aAcceptFailed ; "accept() failed"
.text:08048545      call    puts
.text:0804854A      mov     [esp+2760h+fd], 1
.text:08048551      call    exit
.text:08048556 ; -----|-----
.text:08048556
.text:08048556 loc_8048556:                                ; CODE XREF: main+116↑j
.text:08048556      mov     [esp+2760h+len], 2710h
.text:0804855E      mov     [esp+2760h+addr], 0
.text:08048566      lea     eax, [ebp-2720h]
.text:0804856C      mov     [esp+2760h+fd], eax
.text:0804856F      call    memset
.text:08048574      mov     [esp+2760h+len], 270Fh
.text:0804857C      lea     eax, [ebp-2720h]
.text:08048582      mov     [esp+2760h+addr], eax
.text:08048586      mov     eax, [ebp-8]
.text:08048589      mov     [esp+2760h+fd], eax
.text:0804858C      call    read
.text:08048591      mov     [ebp-10h], eax
.text:08048594      cmp     dword ptr [ebp-10h], 0FFFFFFFFh
.text:08048598      jnz     short loc_80485AB
```

# Reversing the Binary (OS X)

- Things to note
  - IDA identifies it as a MACH-O binary
  - Code looks similar to linux binary

# Reversing the Binary (OS X)

```
call    _listen$UNIX2003
mov     [ebp+var_2728], 10h
jmp     short $+2
```

```
loc_1E5F:
lea     eax, [ebp+var_2728]
lea     edx, [ebp+var_2748]
mov     [esp+8], eax
mov     [esp+4], edx
mov     eax, [ebp+var_10]
mov     [esp], eax
call    _accept$UNIX2003
mov     [ebp+var_C], eax
cmp     [ebp+var_C], 1
jnz     short loc_1EA1
```

```
eax, (aAcceptFailed -      ) [ebx] ; "accept() failed"
[esp], eax
_puts
dword ptr [esp], 1
_exit
```

```
loc_1EA1:
lea     eax, [ebp+var_2724]
mov     edx, eax
mov     eax, 2710h
mov     [esp+8], eax
mov     dword ptr [esp+4], 0
mov     [esp], edx
call    _memset
mov     dword ptr [esp+8], 270Fh
lea     eax, [ebp+var_2724]
mov     [esp+4], eax
mov     eax, [ebp+var_C]
mov     [esp], eax
call    _read
mov     [ebp+var_14], eax
cmp     [ebp+var_14], 0FFFFFFFh
jnz     short loc_1EFB
```

# Reversing the Binary (Windows)

- Things to note
  - Uses threads
  - Symbol support (type information)

# Reversing the Binary (Windows)



```
; Attributes: bp-based frame

; int __cdecl main(int argc, const char **argv, const char **envp)
_main proc near

var_4= dword ptr -4
argc= dword ptr 8
argv= dword ptr 0Ch
envp= dword ptr 10h

push    ebp
mov     ebp, esp
push    ecx
call    sub_401A90
push    0                ; lpThreadId
push    0                ; dwCreationFlags
push    0                ; lpParameter
push    offset StartAddress ; lpStartAddress
push    0                ; dwStackSize
push    0                ; lpThreadAttributes
call    ds:CreateThread
mov     [ebp+var_4], eax
```



# Reversing the Binary (Windows)

```
lea     edx, [ebp+addrlen]
push    edx                ; addrlen
lea     eax, [ebp+addr]
push    eax                ; addr
mov     ecx, [ebp+s]
push    ecx                ; s
call    accept
mov     [ebp+var_273C], eax
push    0                  ; flags
push    2710h              ; len
lea     edx, [ebp+buf]
push    edx                ; buf
mov     eax, [ebp+var_273C]
push    eax                ; s
call    recv
lea     ecx, [ebp+buf]
push    ecx                ; void *
call    sub_401910
add     esp, 4
jmp     short loc_401A34
```

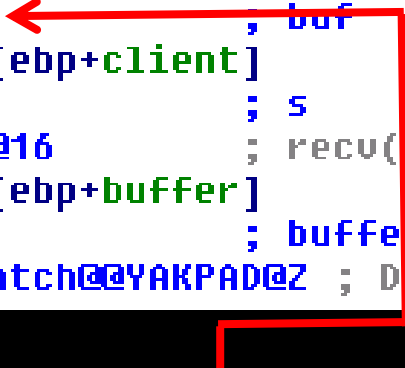


# Reversing the Protocol

- Must identify where network program “recv” network input
  - How the network input is manipulated or understood (protocol parsing)
    - Use cross references to identify network input points

# Reversing the Protocol

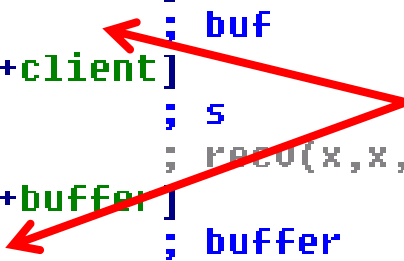
```
push    0                ; flags
push    2710h            ; len
lea     edx, [ebp+buffer]
push    edx              ; buf
mov     eax, [ebp+client]
push    eax              ; s
call    _recv@16         ; recv(x,x,x,x)
lea     ecx, [ebp+buffer]
push    ecx              ; buffer
call    ?Dispatch@@YAKPAD@Z ; Dispatch(char *)
```



- recv stores result in buffer

# Reversing the Protocol

```
push    0                ; flags
push    2710h            ; len
lea     edx, [ebp+buffer]
push    edx              ; buf
mov     eax, [ebp+client]
push    eax              ; s
call    _recv@16         ; recv(x,x,x,x)
lea     ecx, [ebp+buffer]
push    ecx              ; buffer
call    ?Dispatch@@YAKPAD@Z ; Dispatch(char *)
```



- recv stores result in buffer
- buffer is passed to Dispatch function
  - Dispatch takes only one parameter

# Reversing the Protocol

- Dispatch
  - Copies the first four bytes from the buffer into a separate buffer

```
push    4                ; count
mov     eax, [ebp+buffer]
push    eax              ; src
lea     ecx, [ebp+dispatch]
push    ecx              ; dst
call    _memcpy
```

# Reversing the Protocol

- Dispatch

- Calls strtol on first four bytes of packet

- long int strtol(const char \* str, char \*\* endptr, int base);
    - Parses the C string *str* interpreting its content as an integral number of the specified *base*, which is returned as a long int value.

```
push    10                ; ibase
lea     edx, [ebp+p]
push    edx               ; endptr
lea     eax, [ebp+dispatch]
push    eax               ; nptr
call    _strtol
```

# Reversing the Protocol

- Dispatch
  - Recap:
    - Takes recv'd buffer
    - Parses out first four bytes
    - Converts first four bytes to long int
    - **Then enters a switch of valid integers**

# Reversing the Protocol

- Dispatch
  - Protocol accepts the following integers:
    - 1, 2, 3

```
mov     edx, [ebp+opcode]
mov     [ebp+var_10], edx
cmp     [ebp+var_10], 1
jz      short loc_4099F6
cmp     [ebp+var_10], 2
jz      short loc_409A04
cmp     [ebp+var_10], 3
jz      short loc_409A12
jmp     short loc_409A1E
```

Converted integer

Valid integers

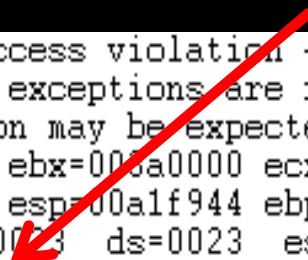
# Opcode: 1

- Where do you want to set a breakpoint?
- Payload:  
[opcode][size][string]
- Judging from disassembly, payload and crash  
what do we suspect the vulnerability to be?



# Crash...

- Where did it crash:



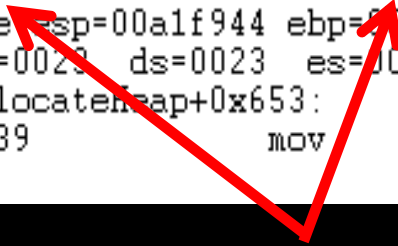
```
(118.88c): Access violation - code c0000005 (first chance)
First chance exceptions are reported before any exception handling.
This exception may be expected and handled.
eax=41414141 ebx=003a0000 ecx=41414141 edx=003a30d0 esi=003a30c8 edi=00000044
eip=7c910f1e esp=00a1f944 ebp=00a1fb64 iopl=0         nv up ei pl zr na pe nc
cs=001b  ss=0023  ds=0023  es=0023  fs=003b  gs=0000             efl=00010246
ntdll!RtlAllocateHeap+0x653:
7c910f1e 8b39                mov     edi,dword ptr [ecx]  ds:0023:41414141=????????
```

- Heap Related?

# Crash...

- Where did it crash:

```
(118.88c): Access violation - code c0000005 (first chance)
First chance exceptions are reported before any exception handling.
This exception may be expected and handled.
eax=41414141 ebx=003a0000 ecx=41414141 edx=003a30d0 esi=003a30c8 edi=00000044
eip=7c910f1e esp=00a1f944 ebp=00a1fb64 iopl=0         nv up ei pl zr na pe nc
cs=001b  ss=0023  ds=0023  es=0023  fs=003b  gs=0000             efl=00010246
ntdll!RtlAllocateHeap+0x653:
7c910f1e 8b39                mov     edi,dword ptr [ecx]  ds:0023:41414141=????????
```



- What do our registers contain?

# Crash...

- Check the call stack for hints

```
ntdll!RtlAllocateHeap+0x653 (FPO: [Non-Fpo])
MSVCR80!_calloc_impl+0x125 (FPO: [Non-Fpo]) (CONV: cdecl) [
MSVCR80!_calloc_crt+0x13 (FPO: [2,0,0]) (CONV: cdecl) [f:\c
MSVCR80!__CRTDLL_INIT+0x1e6 (FPO: [Non-Fpo]) (CONV: cdecl)
MSVCR80!_CRTDLL_INIT+0x1d (FPO: [3,0,0]) (CONV: stdcall) [f
ntdll!LdrpCallInitRoutine+0x14
ntdll!LdrpInitializeThread+0xc0 (FPO: [Non-Fpo])
ntdll!_LdrpInitialize+0x219 (FPO: [Non-Fpo])
ntdll!KiUserApcDispatcher+0x7
```

# Crash...

- Depending on heap state and frequency of heap operations, different crashes can occur

```
(1440.169c): Access violation - code c0000005 (!!! second chance !!!)
eax=00000003 ebx=00000000 ecx=7fffffff edx=00000003 esi=8151d869 edi=81581a33
eip=004012a0 esp=0151d7ec ebp=0151d7f4 iopl=0         nv up ei pl nz na pe nc
cs=001b  ss=0023  ds=0023  es=0023  fs=003b  gs=0000             efl=00010206
*** ERROR: Module load completed but symbols could not be loaded for C:\Users\ap
GreenMan+0x12a0:
004012a0 8a4603          mov     al,byte ptr [esi+3]          ds:0023:8151d86c=??
```

0:001>

# Crash...

- On Windows, we can check the heap integrity...
  - !heap 0 -v

```
##CORRUPTION FOUND at 0x01841A48
    PreviousSize field does not match Size field in previous entry
    Entry->PreviousSize == 0x4141
    PreviousEntry->Size == 0x103
##The above errors were found in segment at 0x01840000
6:  01e60000
    Segment at 01e60000 to 01ea0000 (00001000 bytes committed)
    Flags:                08001002
    ForceFlags:            00000000
    Granularity:           8 bytes
    Segment Reserve:       00100000
    Segment Commit:        00002000
    DeCommit Block Thres:  00000200
    DeCommit Total Thres:  00002000
```

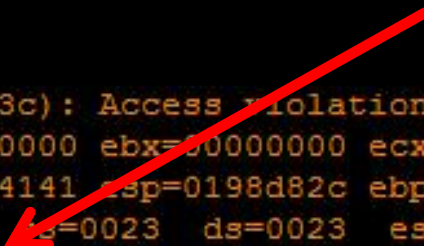
```
0:001> !heap 0 -v
```

# Opcode: 2

- Payload:  
[opcode] [string]
- Judging from disassembly, payload and crash  
what do we suspect the vulnerability to be?

# Crash...

- Where did it crash:



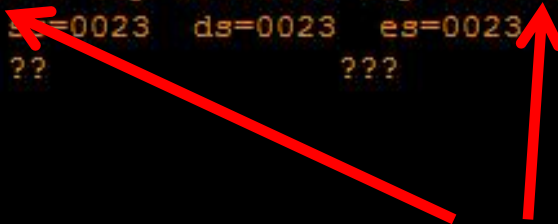
```
(14dc.143c): Access violation - code c0000005 (!!! second chance !!!)
eax=00000000 ebx=00000000 ecx=0198ff74 edx=5f130002 esi=00000000 edi=00000000
eip=41414141 esp=0198d82c ebp=41414141 iopl=0         nv up ei pl zr na pe nc
cs=001b  e8=0023  ds=0023  es=0023  fs=003b  gs=0000             efl=00010246
41414141 ??                ???
```

0:001>

# Crash...

- What do our registers contain?

```
(14dc.143c): Access violation - code c0000005 (!!! second chance !!!)
eax=00000000 ebx=00000000 ecx=0198ff74 edx=5f130002 esi=00000000 edi=00000000
eip=41414141 esp=0198d82c ebp=41414141 iopl=0         nv up ei pl zr na pe nc
cs=001b  eip=0023  ds=0023  es=0023  fs=003b  gs=0000             efl=00010246
41414141 ??                ???
```



0:001>



# Crash...

- What does our call stack look like

```
0:001> kvn
# ChildEBP RetAddr  Args to Child
WARNING: Frame IP not in any known module. Following frames may be wrong.
00 0198d828 41414141 41414141 41414141 41414141 0x41414141
01 0198d82c 41414141 41414141 41414141 41414141 0x41414141
02 0198d830 41414141 41414141 41414141 41414141 0x41414141
03 0198d834 41414141 41414141 41414141 41414141 0x41414141
04 0198d838 41414141 41414141 41414141 41414141 0x41414141
05 0198d83c 41414141 41414141 41414141 41414141 0x41414141
06 0198d840 41414141 41414141 41414141 41414141 0x41414141
07 0198d844 41414141 41414141 41414141 41414141 0x41414141
08 0198d848 41414141 41414141 41414141 41414141 0x41414141
09 0198d84c 41414141 41414141 41414141 41414141 0x41414141
0a 0198d850 41414141 41414141 41414141 41414141 0x41414141
0b 0198d854 41414141 41414141 41414141 41414141 0x41414141
0c 0198d858 41414141 41414141 41414141 41414141 0x41414141
0d 0198d85c 41414141 41414141 41414141 41414141 0x41414141
0e 0198d860 41414141 41414141 41414141 41414141 0x41414141
0f 0198d864 41414141 41414141 41414141 41414141 0x41414141
10 0198d868 41414141 41414141 41414141 41414141 0x41414141
11 0198d86c 41414141 41414141 41414141 41414141 0x41414141
12 0198d870 41414141 41414141 41414141 41414141 0x41414141
13 0198d874 41414141 41414141 41414141 41414141 0x41414141
```

```
0:001>
```

# Opcode: 3

- Payload:  
[opcode] [specific string]
- Judging from disassembly, payload and crash  
what do we suspect the vulnerability to be?

# Conclusion of Session II

**Questions?**

**E-mail the mailing list if you have additional questions**

We are subscribed as well

**Alternatively**

Our gmail usernames are aportnoy and petersilberman

**Thanks!**